

Slide 1

Administrivia

- (None?)

Slide 2

Binary Relations — Review

- Idea of a binary relation is to express relationship between pairs of elements of a set. Formal definition is in terms of sets of ordered pairs.
- Several properties of interest:
 - ρ is *reflexive* if $x \rho x$ for all $x \in S$.
 - ρ is *symmetric* if $(x \rho y) \rightarrow (y \rho x)$ for all $x, y \in S$.
 - ρ is *transitive* if $(x \rho y) \wedge (y \rho z) \rightarrow (x \rho z)$ for all $x, y, z \in S$.
 - ρ is *antisymmetric* if $(x \rho y) \wedge (y \rho x) \rightarrow (x = y)$ for all $x, y \in S$.

Partial Ordering

Slide 3

- Idea: Generalize idea of “ordering” to include relations where not all pairs of elements can be ordered.
- Definition: ρ is a partial ordering if it's reflexive, antisymmetric, and transitive.
- Examples: \leq on integers or reals, \subseteq on sets.
- If finite, can represent with “Hasse diagram” (see examples in textbook).
- Related terms:
 - Successor, predecessor, immediate predecessor, immediate successor.
 - Least, greatest elements.
 - Minimal, maximal elements.

Equivalence Relation

Slide 4

- Idea: Generalize idea of “equals” to include relations where pairs of elements are equivalent but not identical.
- Definition: ρ is an equivalence relation if it's reflexive, symmetric, and transitive.
- Examples: $=$ on integers or reals, $(x \bmod n) = (y \bmod n)$ for some n .
- Related terms/ideas:
 - Equivalence classes.
 - Partition of a set.

Closures

- We can also talk about the “closure” of a relation with respect to a property — the smallest superset of the relation that has the property.
- Example: Define relation ρ on integers such that $x \rho y$ iff $y = x + 1$. What is the transitive closure of ρ ?

Slide 5

Uses of Partial Orderings

- One thing a partial ordering (reflexive, symmetric, transitive relation — think “generalized \leq ”) can express — ordering constraints among tasks.
- We’ll look at one application — topological sorting. PERT charts discussed in book.

Slide 6

Topological Sorting

Slide 7

- Idea here is to take a partial ordering and find a way to extend it to a “total” ordering (i.e., add pairs so that for every x and y either $x \rho y$ or $y \rho x$. How is this useful? e.g., find a way to “schedule” interdependent tasks.
- Notice that there could be more than one way to do this for a given partial ordering.
- How to do this? Next slide . . . (May not be covered in class.)

Topological Sorting, Continued

Slide 8

- Algorithm for finding a way to extend a partial ordering — “topological sort”:
- Start with set S and partial ordering ρ on S . Idea is to turn S into a sequence x_1, x_2, \dots such that $(x_i \rho x_j) \rightarrow (i \leq j)$.
- The algorithm might look like this in pseudocode:
 - while (S not empty)
 - pick a minimal element x in S
 - make it the next element of the sequence and remove it from S
 - end while
- Does this work? i.e., does it produce an ordering that extends ρ ? True if we can be sure that for x and y with $x \rho y$ x is picked before y .

Minute Essay

- None — quiz.

Slide 9