

### Administrivia

- Reminder: Quiz 2 Wednesday.

Slide 1

### Minute Essay From Last Lecture

- Convert 34 (base 10) to binary (base 2).
- Convert 34 (base 10) to hexadecimal (base 16).
- Convert 19 (base 16) to decimal (base 10).
- Convert -34 (base 10) to binary, using 16-bit two's complement notation.

Slide 2

## Arrays and Pointers

- Recall (or observe) that in C (and C++) arrays and pointers are “the same”:  
`void foo(char msg[])` same as `void foo(char *msg)`.

- Consider two ways of setting all elements of an array to 0:

```
void clear1(int a[], int n) {
    for (int i = 0; i < n; ++i)
        a[i] = 0;
}
```

```
void clear2(int *a, int n) {
    for (int *p = a; p < a+n, ++p)
        *p = 0;
}
```

Once upon a time, people interested in writing fast code were told to use `clear2`. Why?

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## Arrays and Pointers, Continued

- Compare code for `clear1` (left) and `clear2` (right):

<pre> add \$t0, \$zero, \$zero # i loop: add \$t1, \$t0, \$t0 # 2*i       add \$t1, \$t1, \$t1 # 4*i       add \$t2, \$a0, \$t1 # addr(a[i])       sw \$zero, 0(\$t2)       addi \$t0, \$t0, 1       slt \$t3, \$t0, \$a1       bne \$t3, \$zero, loop </pre>	<pre> add \$t0, \$a0, \$zero # p loop: sw \$zero, 0(\$t0)       addi \$t0, \$t0, 4       add \$t1, \$a1, \$a1 # 2*n (*)       add \$t1, \$t1, \$t1 # 4*n (*)       add \$t2, \$a0, \$t1 # a+n (*)       slt \$t3, \$t0, \$t2       bne \$t3, \$zero, loop </pre>
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- Which is faster? (Look at the instructions marked with \*.)

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### Binary Versus Decimal

- In decimal (base 10) notation, each digit is multiplied by a power of 10. Same idea for binary (base 2), but using powers of 2.
- So, converting from binary to decimal is easy (if tedious), working from definition. Example?
- Converting from decimal to binary? Repeatedly divide by 2 and record remainders ...

We could describe this as a recursive algorithm for computing  $bits(n)$ :

- Base case is  $n < 2$ ; trivial.
- For recursive step, divide  $n$  by 2 to get quotient  $q$  and remainder  $r$ . Then  $n = 2q + r$ , and:
  - The last bit of  $bits(n)$  should be  $r$ .
  - The remaining bits are  $bits(q)$ , left-shifted by 1.
- Terminology: “Least significant” and “most significant” bits.

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### Binary Versus Hexadecimal

- Binary is useful for showing real internal state but not very compact. Decimal is compact but not so easy to convert to/from binary.
- We might notice — easy to convert to/from a base that’s a power of 2. Hence the use of “octal” (base 8) and “hexadecimal” (base 16). For the latter, we need more than 10 digits, so we use “A” through “F”.

Examples?
- Notice that we can also convert directly to/from decimal, much as we did for binary.

Slide 7

## Representing Integers

- Representing non-negative integers is easy — convert to binary and pad on the left with zeros.
- What about negative integers?
- Could try using one bit for sign, but then you have +0 and -0, and there are other complications.
- Or . . . consider a car odometer — in effect, representable numbers form a circle, since adding 1 to largest number yields 0.

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## Representing Integers, Continued

- We could implement the car-odometer idea in binary, and then choose where to “cut the circle” (between smallest and largest):
  - Between 0 and all ones — unsigned integers.
  - Between largest number with 0 as the MSB and smallest number with 1 as MSB — “two’s complement” signed integers.
- Notice that with the two’s complement scheme, +1/-1 moves us “around the circle” — nothing special needed for negative numbers.
- Notice that if we have  $n$  bits, adding  $2^n$  to  $x$  gives us  $x$  again. This leads to an easy way to compute  $-x$ : Compute  $2^n - x$ , and notice that
 
$$2^n - x = (2^n - 1) - x + 1$$
 which is very easy to compute . . .  
 Examples?

### Minute Essay

- Convert  $30_{10}$  to binary and then to hexadecimal.
- Convert  $-30_{10}$  to 16-bit two's complement notation; show in binary and hexadecimal.
- Convert  $2a_{16}$  to decimal.

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