

Sorted Binary Trees (Binary Search Trees) — Recap
Key property — everything in the left subtree is smaller than the root, and everything in the right is bigger.
Why is this useful? If you want a data structure to hold a collection that will be searched frequently, what are the choices? this one may be better than the others.
Sample programs page has example code (and also code for a recursive implementation of a sorted linked list).





## Heaps — Inserting and Removing Elements

• Basic idea is that we start an insert operation, or a remove-the-largest operation, in a way that maintains the property that the tree is complete. (So, for inserting we add to the lowest level of the tree at the right, and for removing the largest element we replace the top element with the rightmost node at the lowest level.)

Slide 5

- But that breaks the other property (each node larger than its children), so  $\ldots$ 
  - For inserting, we "walk" from the new node up the tree, exchanging with the parent node if they're not in the right order (as sketched in class).
  - For removing, we "walk" from the replaced root node down the tree, exchanging with the largest child if at least one child is larger (also as sketched in class).



- How to store a complete tree in an array? Store "levels" from top to bottom, each level left to right, as sketched in class.
- How to navigate up and down the tree? figure out how to map from a node's index to the indices of its parent and left and right children ...

- Left child of node n is 2n + 1, right child is 2n + 2.
- Parent of node n is (n-1)/2, where the division is integer division (which discards the remainder).







Minute Essay
Sketch what a sorted binary tree of integers would look like after adding the following:
5, 4, -1, 10, 6, 20.
Now sketch what a heap of integers (ordered to put smallest values at the top) would look like after adding the same values.

