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Administrivia

- Reminder: Homework 3 due today.

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Recursion and Recursive Definitions — Review/Recap

- Idea of recursion closely related to idea of induction — “build on previous smaller cases”.
- First look at recursive definitions. To define something recursively:
 - Define one or more “base cases”.
 - Define remaining cases in terms of other (“smaller”) cases.
- Last time we looked at recursive definitions of sequences and sets. (Notice revision to slide about sets.)

Recursive Definitions — Operations

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- Example — factorial.
- Example — multiplication of non-negative integers, defined in terms of addition.
- Example — (integer) division of a non-negative integer by a positive integer, defined in terms of subtraction.

Recursive Algorithms

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- Recursive definitions of sequences or operations often can be turned into recursive algorithms with little effort.
- Examples — function to compute n -th Fibonacci number, function to do division by repeated subtraction.
- Efficiency considerations:
 - In terms of computer time/memory usage, recursion is almost always worse than iteration — but not always, and sometimes not much worse.
 - In terms of human effort to get program running correctly, recursion may be much better.
- Examples in text — selection sort and binary search. Quicksort and mergesort are other good ones.

Reasoning About Recursive Algorithms

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- A recursive algorithm “works” if:
 - It works for the base case(s).
 - For other cases, it works *assuming* the recursive calls work.
 - The recursion eventually stops — recursive calls are always “smaller”, and eventually reduce to base cases.
- We could formalize this as a proof by induction.

Recurrence Relations

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- Recall the silly example of defining a sequence recursively:

$$S(1) = 1$$

$$S(n) = S(n - 1) \times 10, \text{ for } n > 1$$

Expanding out some terms, it seems fairly obvious that an equivalent definition would be $S(n) = 10^{n-1}$ — a “closed-form solution” to the recurrence relation given in the second line of the definition.

- We’ll look at various ways to get from a recursive definition to a closed-form one, because the latter are easier to compute, but sometimes it will be much easier to write down the definition recursively.

Solving Recurrence Relations, Continued

- For the silly example

$$S(1) = 1$$

$$S(n) = S(n-1) \times 10, \text{ for } n > 1$$

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we guessed a solution of $S(n) = 10^{n-1}$. Can we verify that this is the same as the recursive definition? yes, via a proof by induction . . .

- Call this method “expand, guess, verify”.
- Try another example — section 2.5 problem 3.

Solving Recurrence Relations, Continued

- Is there another way? In general, probably not, but there are some formulas for some frequently-occurring special cases.
- One is “first-order linear” recurrence relations with constant coefficients. If

$$S(n) = cS(n-1) + g(n)$$

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then we can show (see textbook for derivation) that

$$S(n) = c^{n-1}S(1) + \sum_{i=2}^n (c^{n-i}g(i))$$

- Apply this to the two problems we did earlier — we should get the same results.

Minute Essay

- Consider the following recursive definition of a sequence:

$$S(1) = 1$$

$$S(n) = 10S(n-1) + 1, \text{ for } n > 1$$

What are $S(1), S(2), \dots, S(5)$?

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Minute Essay Answer

- The first few terms:

$$S(1) = 1$$

$$S(2) = 11$$

$$S(3) = 111$$

$$S(4) = 1111$$

$$S(5) = 11111$$

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