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## Solving Recurrence Relations — Review/Recap

- Goal is to take a recursive definition of a sequence and come up with a "closed-form" (non-recursive) definition of the same sequence.
- One method is what textbook calls "expand, guess, verify".
- No other general method (that I know of!), but formulas for some common special cases. We looked at one last time — if the recursive part of the definition is

$$S(n) = cS(n-1) + g(n)$$

the solution is

$$S(n) = c^{n-1}S(1) + \sum_{i=2}^{n} (c^{n-i}g(i))$$

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## • First find roots $r_1$ and $r_2$ of $t^2 - c_1 t - c_2 = 0$ • and then (if $r_1 \neq r_2$ ) find p and q to solve p + q = S(1) $pr_1 + qr_2 = S(2)$ • and then the formula is $S(n) = pr_1^{n-1} + qr_2^{n-1}$

## Yet Another Special Case

• One more case for which there's a formula is one of interest in analysis of algorithms, especially those that take a "divide and conquer" approach (e.g., quicksort, mergesort, binary search). In math terms, the recursive part is

$$S(n) = cS(n/2) + g(n), \text{ for } n = 2^m, n > 1$$

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• For problems that fit this case, the "expand, guess, verify" method produces the following:

$$S(n) = c^{\log n} S(1) + \sum_{i=\log n}^{n} (c^{\log n - i} g(2^i))$$

## Examples

- Word problem: problem #18 in section 2.5 of textbook.
- Second-order: Fibonacci sequence (problems 24 in 2.4, 31 in 2.5).
- Divide-and-conquer: practice #25 in textbook. (Next time.)

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