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## Memory Hierarchy — Overview Significant overlap between Chapter 5 and material covered in operating-systems course (as I teach it anyway). So, executive-level summary now. A key idea (borrowed from one writer of o/s textbooks): In a perfect world, we could have as much memory as we wanted, and it would be very fast and very cheap. In the real world, there are tradeoffs (e.g., fast versus cheap, fast versus large).



# Memory Hierachy and Caching To exploit temporal locality, can use "caching" — keep copies of frequently-used data in faster but smaller memory. Can do this on multiple levels. To exploit spatial locality, can move data between levels in blocks. Terminology — cache hit, cache miss, cache block/line. Notice that while impact of caching on performance can be significant, it should not affect results (which is why it makes some sense to just ignore it initially).

### Caches (Between Processor and RAM) — Executive-Level Summary

• To make these work, we need:

Some way to map memory address to cache location — can be simple ("direct map") or not.

Some way to say, for each cache location, what memory address it's currently associated with, and whether the data is valid.

- Read "from memory" tries cache first, and then if not found there goes to RAM and updates cache.
- Write "to memory" is maybe more interesting writes to cache but then must at some point write to RAM also — maybe right away (easier to get right but can be slow) or later.

### Virtual Memory — Executive-Level Summary

- Basic idea here is to fake having more RAM than you really have, by keeping some data that would be in RAM on disk. In a sense, RAM is a cache for the "real" memory, on disk(!).
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- Also provides a nice way to support multitasking notion of "processes", each with its own "address space", with an operating system that maps this abstraction onto the hardware, by mapping "program addresses" (in a process's address space) to "physical addresses" (in RAM). Lots of details here, but the basic idea is fairly simple. One big advantage is more control over what data each process can access.



### Caches and Applications Programming Mostly the memory hierarchy (including virtual memory) is managed transparently by a combination of hardware and (operating-system) software, so the first approximation presented in introductory courses (memory is essentially a really big array of bytes, with addresses as indices) is okay, especially if you just want right answers. However, effects on performance can be significant, so if you want right

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• However, effects on performance can be significant, so if you want right answers fast . . .

For single-threaded programs, key idea is to maximize locality (temporal and spatial). Rearranging order in which data is accessed can have a big effect. For multi-threaded programs, also need to consider whether multiple threads need to share access to the same data (problem for correctness too!) or even nearby data ("false sharing" — no effect on correctness but can be slow).

