





- Recall basic ideas (divide address spaces and memory into fixed-size chunks, optionally-but-usually use disk to hold what we hope are less-used parts of processes' address spaces).
- One key issue in making this all work acceptably is how we choose which pages to keep in memory (page replacement algorithm).
- A few more things to consider ...









Sharing Pages and fork

- Duplicating pages is easy but inefficient, especially if the child process is going to call execve or something similar right away. Some systems use "copy-on-write" to improve efficiency.
- Why did the people who designed UNIX require this duplication ... Possibly because it makes some things easy (such as setting up parent/child pipes) and wasn't very costly when designed. Windows' system call for creating processes takes a different approach. Maybe that's better!



Shared Libraries

 One attraction is somewhat obvious — if code for library functions (e.g., printf) is statically linked into every program that uses it, programs need more memory — seems wasteful if processes can share one copy of code in memory.

- Another attraction is that library code can be updated independently of programs that use it. (Is there a downside to that?)
- How to make this happen ... At link time, programs get "stub" versions of functions. References to real versions resolved at load time. Does this remind you of anything? and suggest a possible problem? how to fix?



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Memory-Mapped File I/O Worth mentioning here that some systems also provide a mechanism (e.g., via system calls) to allow reading/writing whole files into/from memory. If there's enough memory, this could improve performance. Example of how this works in Linux — man page for mmap.

