

- Reminder: Reading quizzes 3 and 4 due this week, Homeworks 3a and 3b Monday.
- For the midterm I'm planning a take-home exam, open book/notes like all my exams. (I haven't decided about timed versus untimed.) I'd like to make that available middle of next week and give you about a week to work on it. (More in the minute essay.)

Slide 2

O/S Versus Application Programs — Recap/Review Should seem reasonable to make distinction between what O/S can do and what application programs can do. But how to enforce that? i.e., how to make it as difficult as possible for buggy or malicious application programs to do what they shouldn't? Can this problem be solved completely by clever programming? Consider that most current systems can be asked to load and execute machine-level application code ...

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O/S Versus Application Programs, Continued

- At hardware level, then, need to keep track of which mode we're in and use that information to allow/disallow certain operations (and maybe memory accesses though that could be a separate problem/solution).
- To do this efficiently single bit in a register somewhere, probably a special-purpose one, checked by "privileged" instructions.
- What happens if unprivileged program tries ...? Hardware version of exception — interrupt.
- How to set this bit? privileged operation, or no?





O/S Versus Application Programs, Continued

- So maybe we need memory protection too? but we probably needed that anyway.
- How to make memory protection work? more about that later, but for now again, seems like the only way to do this reliably and efficiently is with help from hardware.
- Most (many?) schemes for memory protection involve some special-purposes registers. Access to these registers privileged mode or not?

Slide 6





Homework 1 Programming Problem, Revisited

• The simple shell you wrote in this assignment created a new process for each command, using fork(), which creates *a full copy of the calling process*, including its program counter, with the intent of using this process to run the desired command. So now you have two processes, a "parent" and a "child" ...

Slide 9

- The parent process should then wait for the child to complete (successfully or not) and then continue with the next command.
- Meanwhile, the child process should use execve to what? If it succeeds, it discards the running program (a copy of the parent process) and executes the program from the specified file, terminating when it's done. What if it doesn't succeed? *The existing program keeps running.* "Oops"? Does this explain behavior that were you puzzled?



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